

PostgreSQL Conference East 2009

# The Art of Indexes

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# Introduction to Indexes

- All About Performance, avoid this:

QUERY PLAN

---

**Seq Scan** on table (cost=0.00..16421.58 rows=1 width=4)

- Each index can dramatically improve performance, both read and write
- But maintaining indexes requires a performance hit



# The Art of Indexes

- Index Type Overview
- Reading EXPLAIN output
- Multi-column indexes
- Expression-based (functional) indexes
- Partial indexes
- Clustered indexes




# The General Syntax

```
CREATE INDEX index_name  
ON table_name ( column1, column2 );
```



# The Full Syntax

```
CREATE [UNIQUE] INDEX [CONCURRENTLY]
  index_name ON [schema.]table_name
  [USING index_type]
  ( column1 or (expression1), ... )
  [WITH (FILLFACTOR = value)]
  [TABLESPACE tablespace_name]
  [WHERE predicate]
;
```



# Reading EXPLAIN

- EXPLAIN shows the planner output
- `<query> → EXPLAIN <query>`
- EXPLAIN ANALYZE runs query, shows run time
  - INSERT, UPDATE, DELETE -will- change the DB
  - BEGIN;  
EXPLAIN ANALYZE UPDATE ...;  
ROLLBACK;



# Reading EXPLAIN

```
=# EXPLAIN SELECT * FROM orders WHERE customerid = 1111;
```

QUERY PLAN

---

```
Seq Scan on orders (cost=0.00..250.00 rows=2 width=36)  
  Filter: (customerid = 1111)  
(2 rows)
```


```
=# CREATE INDEX orders_cid_idx ON orders (customerid);  
CREATE INDEX
```

```
=# EXPLAIN SELECT * FROM orders WHERE customerid = 1111;
```

QUERY PLAN

---

```
Index Scan using orders_cid_idx on orders  
  (cost=0.00..8.27 rows=2 width=36)  
  Index Cond: (customerid = 1111)  
(2 rows)
```





# Reading EXPLAIN

- A larger example...

```
=# EXPLAIN SELECT * FROM customers INNER JOIN orders USING
(customerid) INNER JOIN orderlines USING (orderid) INNER
JOIN products USING (prod_id) INNER JOIN categories
USING (category);
```

QUERY PLAN

```
-----
Hash Join (cost=2811.58..8440.64 rows=60350 width=479)
  Hash Cond: (products.category = categories.category)
  -> Hash Join (cost=2810.22..7609.47 rows=60350 width=361)
    Hash Cond: (orderlines.prod_id = products.prod_id)
    -> Hash Join (cost=2484.22..5774.72 rows=60350 width=314)
      Hash Cond: (orderlines.orderid = orders.orderid)
      -> Seq Scan on orderlines (cost=0.00..958.50 rows=60350 width=18)
      -> Hash (cost=1859.22..1859.22 rows=12000 width=300)
        -> Merge Join (cost=0.10..1859.22 rows=12000 width=300)
          Merge Cond: (customers.customerid = orders.customerid)
          -> Index Scan using customers_pkey on customers (cost=0.00..963.25 rows=20000 width=268)
          -> Index Scan using orders_cid_idx on orders (cost=0.00..696.22 rows=12000 width=36)
        -> Hash (cost=201.00..201.00 rows=10000 width=51)
          -> Seq Scan on products (cost=0.00..201.00 rows=10000 width=51)
    -> Hash (cost=1.16..1.16 rows=16 width=122)
      -> Seq Scan on categories (cost=0.00..1.16 rows=16 width=122)
(16 rows)
```



# Index Type Overview

- Direct Types:
  - B-tree
  - Hash
- Container Types:
  - GIN
  - GiST

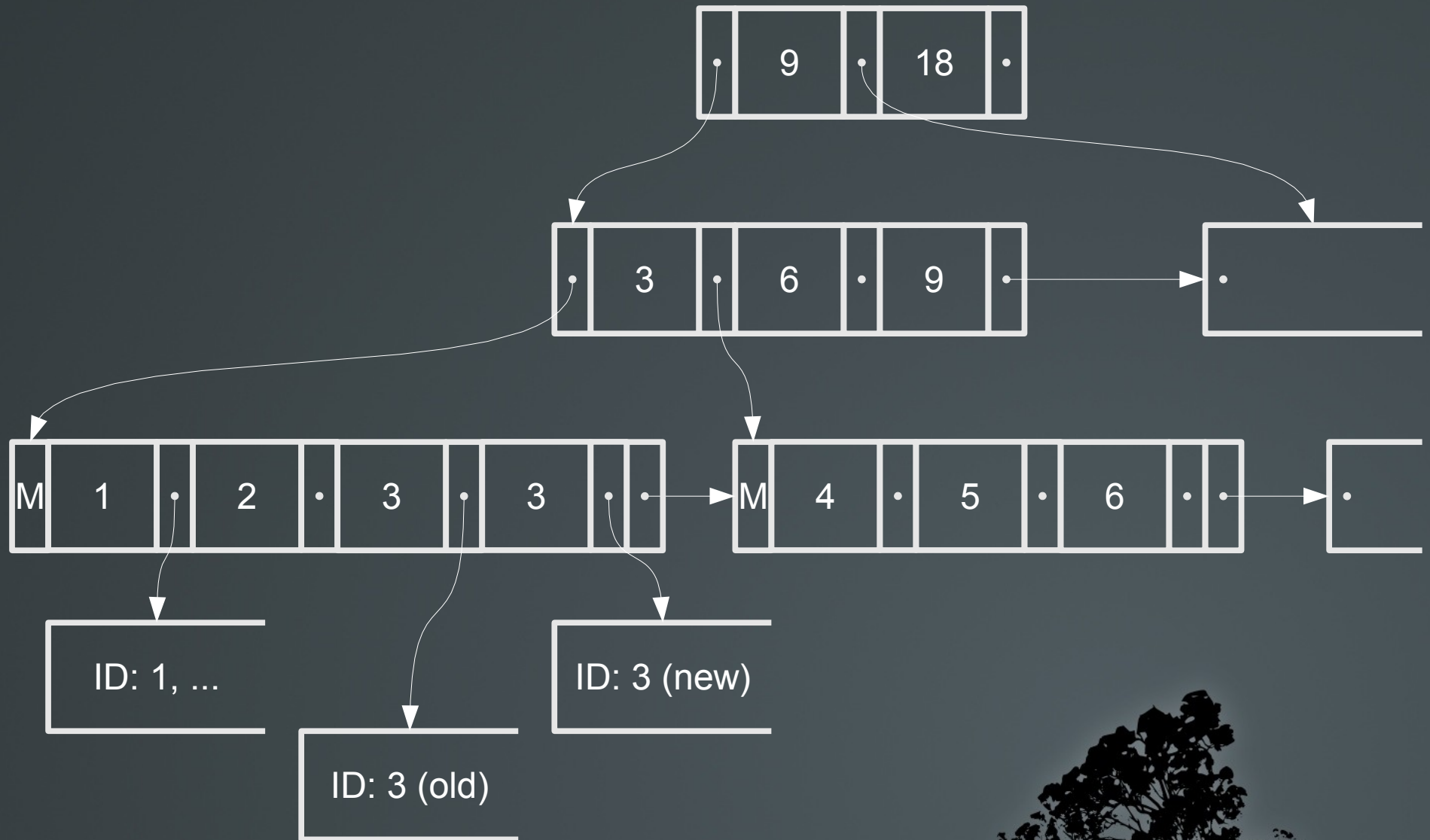


# Index Types: B-tree

- Default index type
- Strengths
  - Supports range queries on common data types (number & string types)
- Weaknesses
  - Upper limit on field size (~1/3 page size)

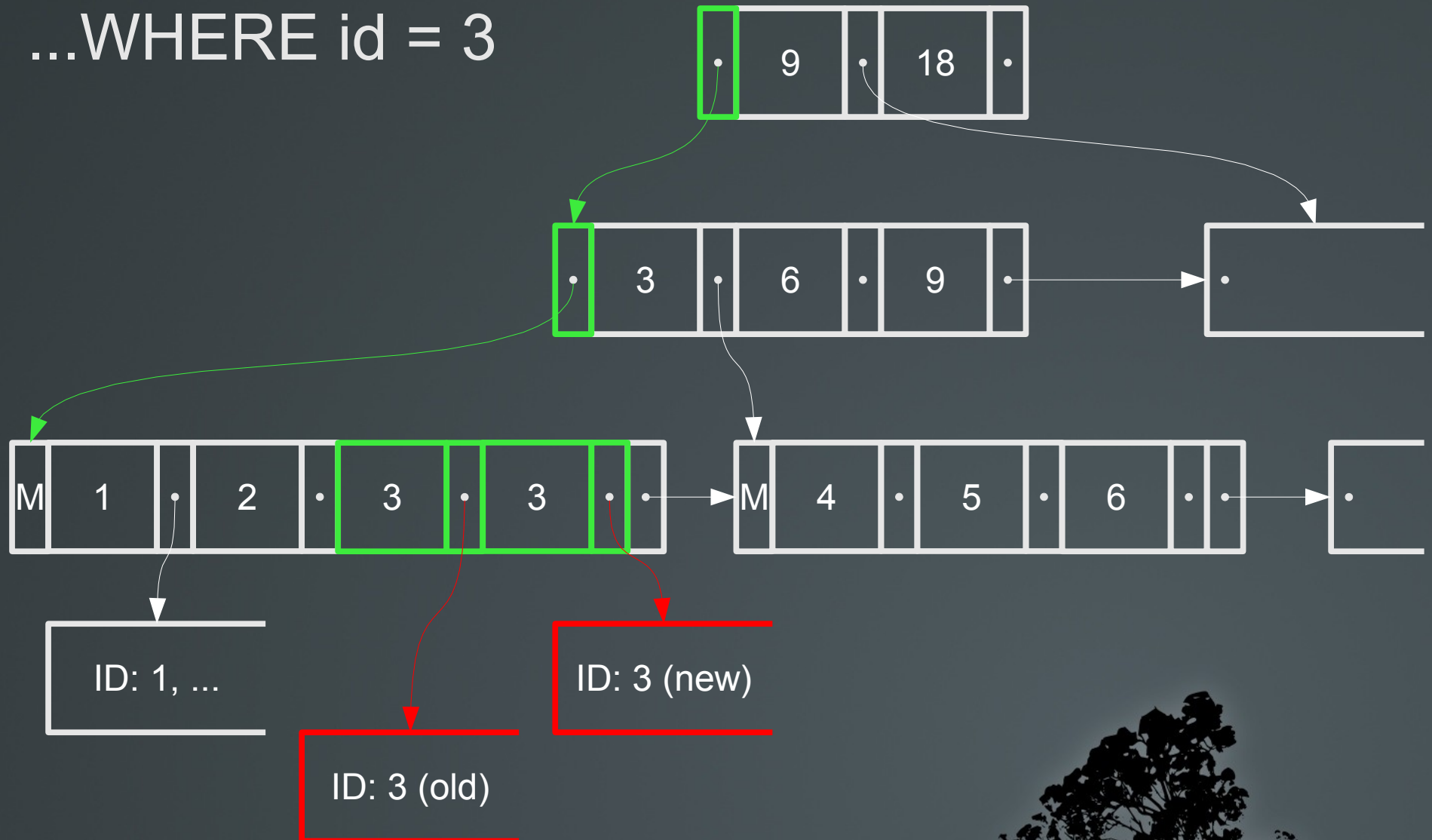


# Index Types: B-tree



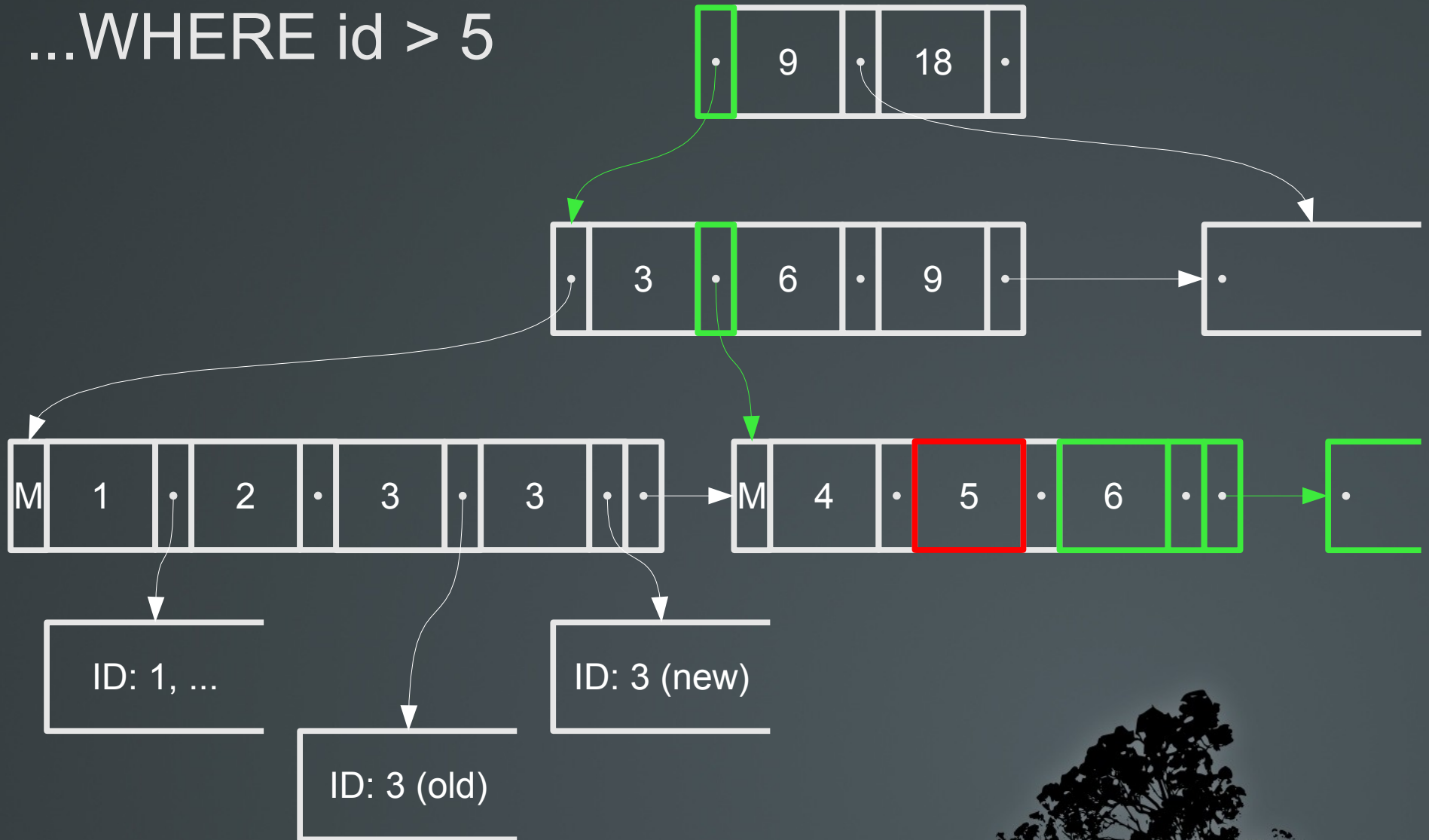
# Index Types: B-tree

...WHERE id = 3



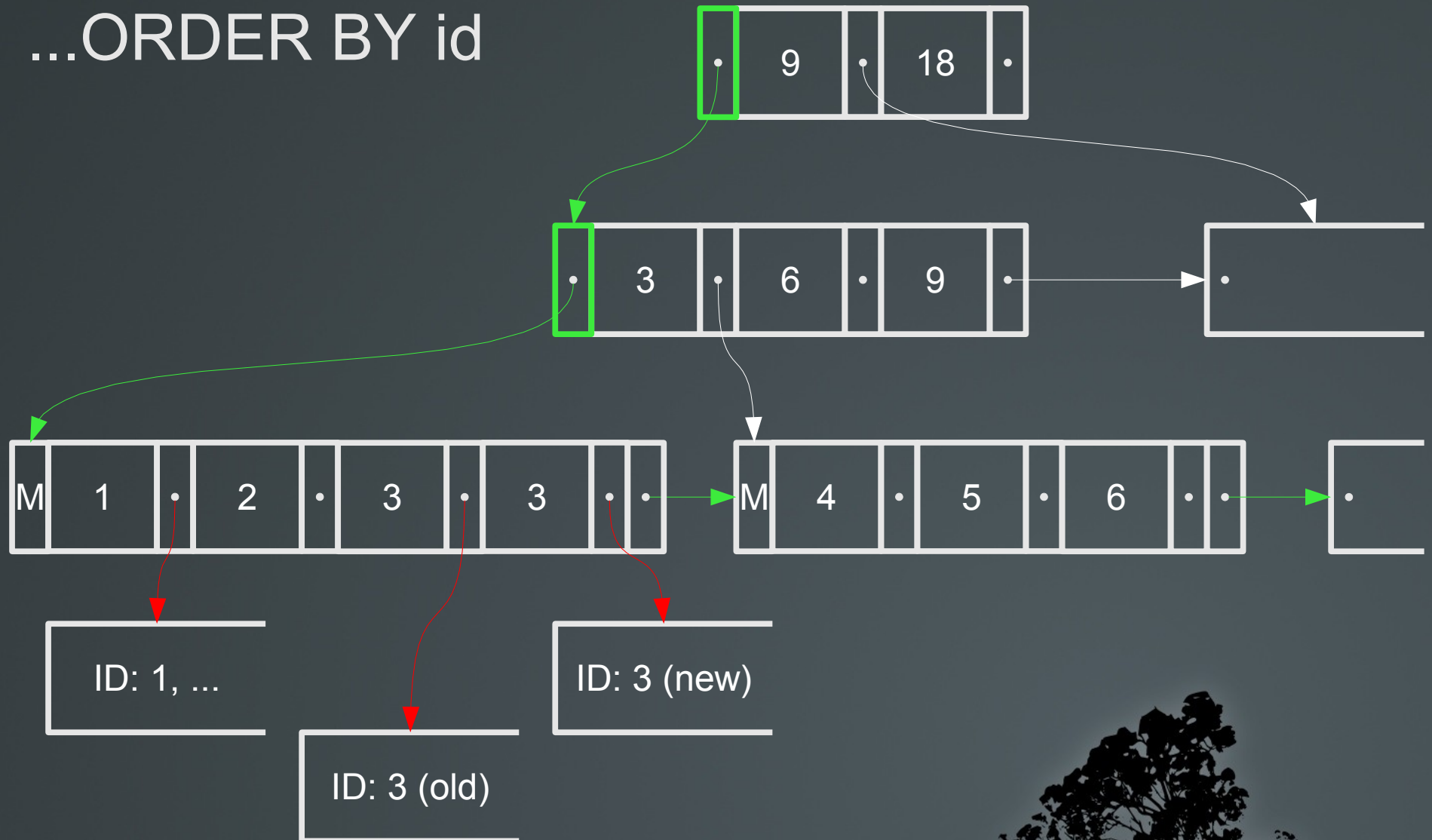
# Index Types: B-tree

...WHERE id > 5



# Index Types: B-tree

...ORDER BY id



# Index Types: Hash

- Strengths
  - Better at handling low-cardinality situations than B-tree
  - Received some improvements in 8.4 (Robert said so.)
- Weaknesses
  - Poor build time, search time about the same as B-tree
  - Only supports equality operator, No mechanism to reduce size, No CLUSTER support, No WAL-logging
- Nah, better off not considering this one





# Index Types: GIN

- Generalized Inverted Index
- Useful for array types (including full text columns)
- Similar to hash indexes...
  - Sets of key/'element' and 'posting list'/matching row pairs
  - Multiple rows per key, Multiple keys per row
- Extensible: Users can define functions for their data types to interface with the index
- B-tree (keys), pointing to a B-tree (or list) of tuples



# Index Types: GiST

- Generalized Search Tree
- Useful for spatial data types (and full text columns)
- Extensible: Provides a general framework to create tree structures based on user functions
- The output of the user functions define how the tree is created (and searched.)



# Index Types: GIN

- Strengths
  - Indexing columns with multiple keys
  - 8.4 will support multiple columns
- Weaknesses
  - Only supports equality (type) operators
  - ... Though 8.4 supports partial matches
  - Can be expensive to construct



# Index Types: GiST

- Strengths
  - Very versatile, if you have a custom data type
  - Useful for geometric types (and PostGIS!)
- Weaknesses
  - Tend to be slower to search
  - Not as useful for built-in data types



# Full Text Index Strategy

- Both GiST and GIN support full text searches.....
- Which one to use is a matter of trade offs:
  - GiST is always "lossy" and prone to false index matches
  - But it's fast to maintain!
  - GIN is big and slow, and has an operator quirk
  - But it's faster to query, and supports larger documents!
- The Fine Manual's rule of thumb...
  - Static data (e-books) GIN; Dynamic data (forums) GiST.



# Invoking Your New Index

- Run `ANALYZE` on the table – update statistics
- If it still doesn't work, force it to use the index...
  - `SET enable_seqscan = off;`
  - `SET enable_nestloop = off;`
  - `EXPLAIN ANALYZE` (or `\timing` in `psql`)
- If it *still* doesn't work, make sure the index definition (including selected expressions or columns, partial index predicate) matches.




# Indexing Strategy Decisions

- Simple indexing tricks: Concurrent builds, UNIQUE
- Multiple columns, versus many single-column
- Indexes on Expressions / Functions
- Partial Indexes
- Clustering Indexes



# Create Index CONCURRENTLY

- The index build process: Expensive
  - `CREATE INDEX` obtains a table-wide `SHARE` lock (i.e. all write activity is blocked.)
  - `CONCURRENTLY` allows write activity, naturally
    - Performs two scans over the table (separate transactions)
    - Has to wait for previous transactions to finish
    - If it runs into problems (unique, function errors) ... invalid index just hangs around out there
    - `DROP` and try again...
- 



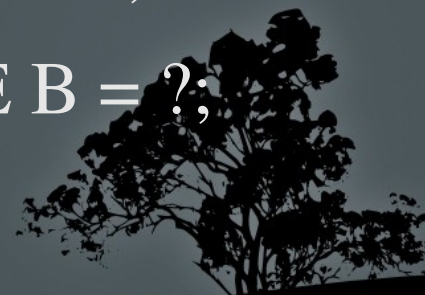
# The UNIQUE Index

- (Primary Key) Unique Constraint ==> Unique Index
- Index is checked for existing values
- Remember: NULL != NULL
- (Primary Key) Multiple columns are supported
- Usually considered to be an implementation detail (in other words, use `ADD CONSTRAINT UNIQUE`)  
... but there's a couple cool things you can do with it.



# Multi-Column Indexes

- Supported by B-tree, GiST, and GIN (8.4)
- Most useful in "parent/child" column relationships
  - Major, minor device numbers
  - Order number, item number in an order detail table
- An index on columns (A, B)...
  - Really useful for `SELECT ... WHERE A = ? AND B = ?;`
  - Still useful for `SELECT ... WHERE A = ?;`
  - Not so useful for `SELECT ... WHERE B = ?;`



# Multiple Single-Column Indexes

- Bitmap Index Scan: use multiple indexes together
- Each index produces a bitmap of tuples (or pages)...
- The bitmaps are AND'ed (or OR'ed) together...

Tuple:	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>	<u>7</u>	<u>8</u>	
Idx1:	0	1	0	0	1	0	0	1	
Idx2:	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>	&&
Result:	0	1	0	0	0	0	0	1	



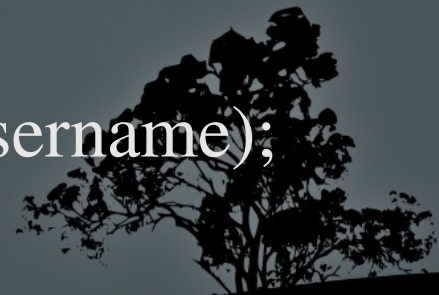
# Multiple Single-Column Indexes

- So multi-column indexes are no longer as important.
- An index on column (A) + an index on column (B)...
  - Useful for `SELECT ... WHERE A = ?;`
  - Useful for `SELECT ... WHERE B = ?;`
  - Useful for `SELECT ... WHERE A = ? AND B = ?;`
- But you have to maintain 2 (or more) indexes



# Expression-Based Indexes

- Indexes can be built on the result of an expression
- Classic example: Usernames
  - `CREATE INDEX ... ON users (lower(username));`
  - `SELECT ... FROM users WHERE lower(username) = '?';`
- Can be expensive to build (and maintain)
- Any functions used must be marked `IMMUTABLE`
- That cool thing we talked about...
  - `CREATE UNIQUE INDEX ... lower(username);`



# The Partial Index

- Restrict what an index will consider
- Similar to `SELECT`'s `WHERE` clause
- `CREATE INDEX ... WHERE balance > 1000;`
  - Allow quick access to 'interesting' data, small index
  - The rest of the table accessed through slower methods
- The other cool thing we talked about...
  - `CREATE UNIQUE INDEX ... (username) WHERE active_account = true;`



# Clustering an Index

- Rewrites a physical table in index order
  - Like `VACUUM FULL ...` but useful!
- A one-time operation, index order is not maintained
  - (But `FILLFACTOR` can help)
- Needs twice the disk space, and an exclusive lock
- Very useful for range queries...
  - `WHERE created BETWEEN '2009-03-01' and '2009-04-01'`
- Run `ANALYZE` afterward, to update statistics



# What Index Options Are Right?

?





# The Art of Indexes

- Index Type Overview
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# The Art of Indexes

Slides...

<http://dbahumor.com/indexes>

